

Programming Languages: Syntax

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Announcements and Reminders

- Sound and visuals check.
- Main online resource: Class Web site:
 - <http://chaw.eip10.org/cos301/>
 - also linked from my Web page, etc.
 - Includes (is) **syllabus**.
- Brightspace for **some things only**.
 - **discussion forum**.
 - homework and other **submissions**.
- **Homework HW01** due soon, not trivial.
- Rules for use of **Gen AI**.

Plan for today

- Homework HW01 Q&A.
- (Continuing) Material from Chapter 2 of the textbook.
 - (Review) *[E]BNF*, *CFG*.
 - Theory of CFGs etc.
 - *Parse trees* and *abstract syntax trees (AST)*.
- (Review) Bigger picture question (related to homework):
 - How to implement a simple language like:
 - $x = 5 + 3$
 - $y = 48 / (4 * 4)$
 - $z = x + 2 * y$
 - etc.

(Review) Terminals and nonterminals

- *terminals* or *tokens*
 - \approx granules of the program source that are not analyzed internally by the CFG
 - but may be analyzed internally by the *lexer*.
 - examples
 - `=`
 - `;`
 - `avonum`
 - `6.022E23`
- *nonterminals* or *syntactic categories*
 - have components that are specified and analyzed by a CFG
 - examples
 - *assignment statement*: `avonum = 6.023;`
 - *return statement*: `return 42;`
 - *predicate*: `avonum > 42`

(Review) BNF: Backus-Naur Form(at)

- BNF spec = *set* of rules
 - N.B.: above spec is in a meta meta language.
- Each rule has the form:
 - *nonterminal* ::= sequence of *terminals* and *nonterminals*
 - again a meta-meta-language spec.
 - ::= = "is" or "is composed of" or "can be replaced by"
- Examples
 - *assignment-statement* ::= *variable-name assignment-operator rval* ; ;
 - meta-language (BNF): ::=, ;
 - language: (C-like): ;
 - *assignment-operator* ::= = ;
 - *statements* ::= *statement statements* | ;
 - | is short-hand for multiple rules with same LHS.

(Review) BNF example

- from textbook

```
<primitive-type> ::= boolean | char | byte | short | int | long | float | ...  
<argument-list> ::= <expression> | <argument-list> , <expression>  
<selection-statement> ::= if ( <expression> ) <statement>  
~~~~~I| if ( <expression> ) <statement> else <statement>  
~~~~~I| switch ( <expression> ) <block>  
~~~~~I;
```

- Exercises

- For each component above: Is it in language or meta-language?
- Describe in English as precisely as possible.
- Provide illustrative examples (in the *language*) making reasonable assumptions.

(Review) EBNF = Extended BNF

- BNF + some convenience features
- foo? or $[\text{foo}]$ = optional foo
 - exercise: language v. metalanguage elements above
- foo^* or $\{\text{foo}\}$ = a sequence of zero or more foo
- foo^+ = a sequence of one or more foo
- parentheses (in metalanguage; language may have them too!)
 - $(\text{foo bar})^+$ = sequence of one or more instances of foo bar
 - $(())^+$ = sequence of one or more $()$.
 - !!

(Review) Discussion from end of previous class

- $(a\ b)^* \stackrel{?}{\equiv} (a\ b)^+?$
- $(a\ b)^+? \stackrel{?}{\equiv} (a\ b)^{?+}$
- etc.
- Elements of language, metalanguage, meta-metalanguage.
- Practice proving and disproving such statements.

CFG formally

- $G = (N, T, P, S)$
 - N : a set of symbols (*nonterminals*)
 - T : another set of symbols (*terminals*)
 - $N \cap T = \emptyset$
 - P : set of *productions*
 - each of form $n \rightarrow \alpha$ where
 - where $n \in N$ and $\alpha \in \{N \cup T\}^*$
 - $S \in N$: special nonterminal called *start*

CFG for infix expressions

- example from the textbook

Derivations

- Main question: Can a given string of tokens (terminals) be generated by a given grammar?
 - If so, how? Show the steps starting with the start symbol.
 - Is the derivation unique?
 - Is the leftmost derivation unique?
- Switch to example and practice in textbook.

Parse Trees

- Informally, a tree that has S as root and the children of each node are the items on the RHS of the rule that was used to replace the corresponding nonterminal. (Leaves correspond to terminals.)
- Example for infix expressions.
- Switch to practice problems.

Abstract Syntax Trees

- closely related to, but different from, parse trees.
- abstract away unimportant details such as order in which a sequence of nonterminals is expanded.
- Two main changes:
 - Nonterminal nodes are replaced by *corresponding* parts of the input sentence.
 - Unit productions are collapsed.
- Example for infix expressions.
- Practice problems in textbook.

Summary

- (E)BNF, CFGs, derivations, parse trees, ASTs.
- Bigger picture:
 - Process source code into AST
 - then we can interpret or compile it etc.
- Class Web site: <http://chaw.eip10.org/cos301/>